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MULTI-DIMENSIONAL SETS RECOGNIZABLE IN ALL ABSTRACT NUMERATION SYSTEMS

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Abstract. We prove that the subsets of $\mathbb{N}^d$ that are S-recognizable for all abstract numeration systems S are exactly the 1-recognizable sets. This generalizes a result of Lecomte and Rigo in the one-dimensional setting.

Mathematics Subject Classification. 68Q45.

1. Introduction

In this paper we characterize the subsets of $\mathbb{N}^d$ that are simultaneously recognizable in all abstract numeration systems (numeration systems that represent a natural number n by the $(n+1)$-th word of a genealogically ordered regular language – see below for the precise definition). Lecomte and Rigo [11] provided such a characterization for the case $d = 1$ based on the well-known correspondence between unary regular languages and ultimately periodic subsets of $\mathbb{N}$. When $d > 1$ we no longer have such a nice correspondence and the situation becomes somewhat more complicated. To obtain our characterization we instead use a classical decomposition theorem due to Eilenberg *et al.* [7]. The motivation for studying such sets comes from the well-known result of Cobham (and its multi-dimensional generalization due to Semenov) concerning the sets recognizable in integer bases.

Let $k \geq 2$ be an integer. A set $X \subseteq \mathbb{N}$ is *k-recognizable* (or *k-automatic*) if the language consisting of the base-k representations of the elements of X is accepted by a finite automaton. A celebrated result of Cobham [5] characterizes the sets that are recognizable in all integer bases $k \geq 2$. Two numbers k and ℓ are *multiplicatively independent* if $k^m = \ell^n$ implies $m = n = 0$. A subset of the integers is *ultimately periodic* if it is a finite union of arithmetic progressions.

Keywords and phrases. Finite automata, numeration systems, recognizable sets of integers, multi-dimensional setting.

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Theorem 1.1 (Cobham)**.** *Let $k, \ell \geq 2$ be two multiplicatively independent integers and let $X \subseteq \mathbb{N}$. The set X is both k-recognizable and ℓ-recognizable if and only if it is ultimately periodic.*

We say that a set $X \subseteq \mathbb{N}$ is 1-*recognizable* if the language $\{a^n : n \in X\}$ consisting of the unary representations of the elements of X is accepted by a finite automaton. It is well-known [6], Proposition V.1.1, that a set is 1-recognizable if and only if it is ultimately periodic.

Lecomte and Rigo [11] introduced the following generalization of the standard integer base numeration systems.

Definition 1.2. An *abstract numeration system* is a triple $S = (L, \Sigma, <)$ where L is an infinite regular language over a totally ordered finite alphabet $(\Sigma, <)$. The map $\mathrm{rep}_S : \mathbb{N} \to L$ is a bijection mapping $n \in \mathbb{N}$ to the $(n+1)$-th word of L ordered genealogically. The inverse map is denoted by $\mathrm{val}_S : L \to \mathbb{N}$. A set $X \subseteq \mathbb{N}$ is S-*recognizable* if the language $\mathrm{rep}_S(X) = \{\mathrm{rep}_S(n) : n \in X\}$ is regular.

Lecomte and Rigo [11] proved that any ultimately periodic set is S-recognizable for any abstract numeration system S. Suppose on the other hand that $X \subseteq \mathbb{N}$ is S-recognizable for every abstract numeration system S. Then in particular, the set X must be 1-recognizable, and hence must be ultimately periodic. We therefore have the following characterization of the sets that are recognizable in all abstract numeration systems.

Theorem 1.3 (Lecomte and Rigo)**.** *A set $X \subseteq \mathbb{N}$ is S-recognizable for all abstract numeration systems S if and only if it is ultimately periodic.*

Rigo and Maes [14] considered S-recognizability in a multi-dimensional setting. This concept was further studied by Charlier *et al.* [4]. For the formal definitions we need to introduce the following "padding" function.

Definition 1.4. If $w_1, \ldots, w_d$ are finite words over the alphabet Σ, the padding map

$$(\cdot)^{\#} : (\Sigma^*)^d \to ((\Sigma \cup \{\#\})^d)^*$$

is defined by

$$(w_1, \ldots, w_d)^{\#} := (w_1 \#^{m-|w_1|}, \ldots, w_d \#^{m-|w_d|})$$

where $m = \max\{|w_1|, \ldots, |w_d|\}$. Here we write (ac, bd) to denote the concatenation $(a, b)(c, d)$.

If $R \subseteq (\Sigma^*)^d$, then

$$R^{\#} = \{(w_1, \ldots, w_d)^{\#} : (w_1, \ldots, w_d) \in R\}.$$

Note that R is not necessarily a language, whereas $R^{\#}$ is; that is, the set R consists of d-tuples of words over Σ, whereas $R^{\#}$ consists of words over the alphabet $(\Sigma \cup \{\#\})^d$.

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Definition 1.5. Let $S = (L, \Sigma, <)$ be an abstract numeration system. Let $X \subseteq \mathbb{N}^d$. The set X is *S-recognizable* (or *S-automatic*) if the language $\mathrm{rep}_S(X)^\#$ is regular, where

$$\mathrm{rep}_S(X) = \{(\mathrm{rep}_S(n_1), \ldots, \mathrm{rep}_S(n_d)) \colon (n_1, \ldots, n_d) \in X\}.$$

Let $k \geq 2$ be an integer. The notions of k-recognizability and 1-recognizability are special cases of S-recognizability. The set X is *k-recognizable* (or *k-automatic*) if it is S-recognizable for the abstract numeration system S built on the language consisting of the base-k representations of the elements of X. The set X is 1-*recognizable* (or 1-*automatic*) if it is S-recognizable for the abstract numeration system S built on a^*.

The multi-dimensional analogue of Cobham's theorem, due to Semenov [16], requires an analogous notion of ultimate periodicity in the multi-dimensional setting.

Definition 1.6. A set $X \subseteq \mathbb{N}^d$ is *linear* if there exists $\mathbf{v}_0, \mathbf{v}_1, \ldots, \mathbf{v}_t \in \mathbb{N}^d$ such that

$$X = \{\mathbf{v}_0 + n_1\mathbf{v}_1 + n_2\mathbf{v}_2 + \ldots + n_t\mathbf{v}_t : n_1, \ldots, n_t \in \mathbb{N}\}.$$

A set $X \subseteq \mathbb{N}^d$ is *semi-linear* if it is a finite union of linear sets.

For more on semi-linear sets see [10]. We can now state the multi-dimensional version of Cobham's Theorem [16].

Theorem 1.7 (Cobham-Semenov)**.** *Let $k, \ell \geq 2$ be two multiplicatively independent integers and let $X \subseteq \mathbb{N}^d$. The set X is both k-recognizable and ℓ-recognizable if and only if it is semi-linear.*

In other words, the semi-linear sets are precisely the sets recognizable in all integer bases $k \geq 2$. One might therefore expect that, as in Theorem 1.3, the semi-linear sets are recognizable in all abstract numeration systems. However, this fails to be the case, as the following example shows.

Example 1.8. The semi-linear set $X = \{n(1, 2) : n \in \mathbb{N}\} = \{(n, 2n) : n \in \mathbb{N}\}$ is not 1-recognizable. Consider the language $\{(a^n\#^n, a^{2n}) : n \in \mathbb{N}\}$, consisting of the unary representations of the elements of X. An easy application of the pumping lemma shows that this is not a regular language.

Observe that in the one-dimensional case, we have the following equivalences: semi-linear $\Leftrightarrow$ ultimately periodic $\Leftrightarrow$ 1-recognizable. However, Example 1.8 shows that these equivalences no longer hold in the multi-dimensional setting. In order to get a multi-dimensional analogue of Theorem 1.3, we must consider the class of 1-recognizable sets, which form a proper subclass of the class of semi-linear sets.

Another well-studied subclass of the class of semi-linear sets is the class of recognizable sets. A subset X of $\mathbb{N}^d$ is *recognizable* if there exists a finite monoid M, a monoid homomorphism $\varphi : \mathbb{N}^d \to M$, and a subset $B \subseteq M$ such that $X = \varphi^{-1}(B)$. When $d = 1$, we have again the following equivalences: recognizable

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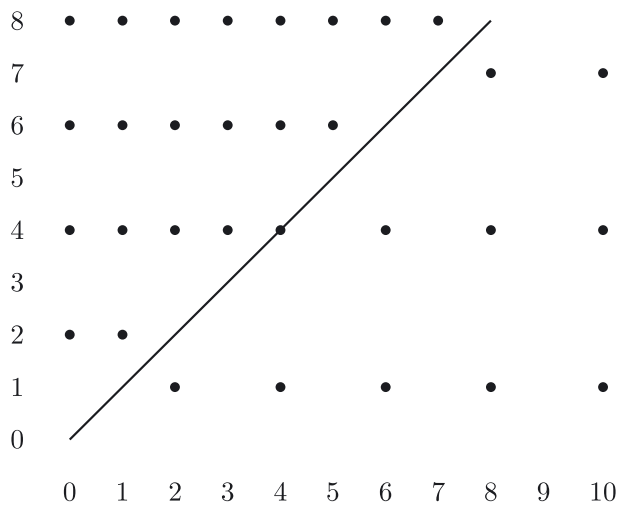


FIGURE 1. The set X of Example 1.10.

$\Leftrightarrow$ ultimately periodic $\Leftrightarrow$ 1-recognizable. However, for $d > 1$ these equivalences no longer hold. An unpublished result of Mezei (see [6], Prop. III.12.2) demonstrates that the recognizable subsets of $\mathbb{N}^2$ are precisely finite unions of sets of the form $Y \times Z$, where Y and Z are ultimately periodic subsets of $\mathbb{N}$. In particular, the *diagonal set* $D = \{(n, n) : n \in \mathbb{N}\}$ is not recognizable [6], Exercise III.12.7. However, the set D is clearly a 1-recognizable subset of $\mathbb{N}^2$. So we see that for $d > 1$, the class of 1-recognizable sets corresponds neither to the class of semi-linear sets, nor to the class of recognizable sets. For further information on recognizable sets, their different characterizations and the classical Cobham-Semenov Theorem, see [3].

Our main result is the following, which generalizes the result of Lecomte and Rigo (Thm. 1.3).

Theorem 1.9. *Let $X \subseteq \mathbb{N}^d$. Then X is S-recognizable for all abstract numeration systems S if and only if X is* 1*-recognizable.*

To illustrate this theorem, we give the following example.

Example 1.10. Let

$$X = \{(2n, 3m+1) : n, m \in \mathbb{N} \text{ and } 2n \geq 3m+1\} \cup \{(n, 2m) : n, m \in \mathbb{N} \text{ and } n < 2m\}$$

(see Fig. 1). It is clear that X is 1-recognizable. Let $S = (L, \Sigma, <)$ be an abstract numeration system. By Theorem 1.3, the sets $\{2n : n \in \mathbb{N}\}$ and $\{3m+1 : m \in \mathbb{N}\}$ are both S-recognizable, and so the set $\{(2n, 3m+1) : n, m \in \mathbb{N}\}$ is also S-recognizable. In other words, the set $\{(\mathrm{rep}_S(2n), \mathrm{rep}_S(3m+1))^{\#} : n, m \in \mathbb{N}\}$ is accepted by a finite automaton. Furthermore, the set $\{(x, y)^{\#} : x, y \in L \text{ and } x \geq y\}$

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is also accepted by a finite automaton, and so by taking the product of these two automata we obtain an automaton accepting

$$\{(\mathrm{rep}_S(2n), \mathrm{rep}_S(3m+1))^{\#} : n, m \in \mathbb{N} \text{ and } 2n \geq 3m+1\}.$$

In the same way we can construct an automaton to accept the set

$$\{(\mathrm{rep}_S(n), \mathrm{rep}_S(2m))^{\#} : n, m \in \mathbb{N} \text{ and } n < 2m\}.$$

Since the union of two regular languages is regular, we see that X is S-recognizable.

2. Proof of our main result

In order to obtain our main result, we will need a classical result of Eilenberg *et al.* [7], Theorem 11.1 (see also [15], Thm. C.1.1). We first need the following definition.

Definition 2.1. Let A be a non-empty subset of $\{1, \ldots, d\}$. Define the subalphabet

$$\Sigma_A = \{x \in (\Sigma \cup \{\#\})^d : \text{the } i\text{-th component of } x \text{ is } \# \text{ exactly when } i \notin A\}.$$

Example 2.2. Let $\Sigma = \{a, b\}$ and $d = 4$. If $A = \{1, 2, 3, 4\}$, then $\Sigma_A = \{(\sigma_1, \sigma_2, \sigma_3, \sigma_4) : \sigma_i \in \Sigma \text{ for } i = 1, 2, 3, 4\}$. If $A = \{2, 3\}$, then $\Sigma_A = \{(\#, \sigma_2, \sigma_3, \#) : \sigma_i \in \Sigma \text{ for } i \in \{2, 3\}\}$. If $A = \{3\}$, then $\Sigma_A = \{(\#, \#, a, \#), (\#, \#, b, \#)\}$.

Theorem 2.3 (Decomposition [7])**.** *Let $R \subseteq (\Sigma^*)^d$. The language $R^{\#} \subseteq ((\Sigma \cup \{\#\})^d)^*$ is regular if and only if it is a finite union of languages of the form*

$$R_0 \ldots R_t, \quad t \in \mathbb{N},$$

where each factor $R_i \subseteq (\Sigma_{A_i})^*$ is regular and $A_t \subseteq \ldots \subseteq A_0 \subseteq \{1, \ldots, d\}$.

Remark 2.4. Theorem 2.3 does not hold if $R^{\#}$ is replaced by an arbitrary language over $(\Sigma \cup \{\#\})^d$. It is only valid due to the definition of the map $(\cdot)^{\#}$.

Example 2.5. Let $R = \{(a^{5n}, a^{6m}) : n, m \in \mathbb{N}\}$. Then $R^{\#}$ is regular, since one can easily construct an automaton that simultaneously checks that the length of the first component of its input is a multiple of 5 and that the length of the second component is a multiple of 6. Moreover, we have

$$R^{\#} = \bigcup_{\ell=0}^{5} (a^{30}, a^{30})^* (a^{5\ell}\#^{\ell}, a^{6\ell})(\#^6, a^6)^* \quad \cup \quad \bigcup_{\ell=0}^{4} (a^{30}, a^{30})^* (a^{5(\ell+1)}, a^{6\ell}\#^{5-\ell})(a^5, \#^5)^*.$$

Observe that each of the languages appearing in the unions above are products of the form described in Theorem 2.3.

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Lemma 2.6. *Let $X \subseteq \mathbb{N}^d$. Then X is 1-recognizable if and only if X is a finite union of sets of the form*

$$\left\{ \sum_{\ell=0}^{t} (c_\ell(n_{\ell,1}, \ldots, n_{\ell,d}) + (b_{\ell,1}, \ldots, b_{\ell,d})) \colon (\forall \ell)(\forall i)\ n_{\ell,i} \in \mathbb{N} \text{ and } \right.$$
$$\left. (\forall \ell)(\forall i)\ (i \notin A_\ell \Rightarrow n_{\ell,i} = 0) \text{ and } (\forall \ell)(\forall i)(\forall j)(i,j \in A_\ell \Rightarrow n_{\ell,i} = n_{\ell,j}) \right\} \quad (2.1)$$

where

- $t \in \mathbb{N}$;
- $A_t \subseteq \ldots \subseteq A_0 \subseteq \{1, \ldots, d\}$;
- $c_0, \ldots, c_t \in \mathbb{N}$;
- $(\forall \ell)(\forall i)\ b_{\ell,i} \in \mathbb{N}$;
- $(\forall \ell)(\forall i)\ (i \notin A_\ell \Rightarrow b_{\ell,i} = 0)$; *and*
- $(\forall \ell)(\forall i)(\forall j)\ (i,j \in A_\ell \Rightarrow b_{\ell,i} = b_{\ell,j})$.

Proof. Let $\Sigma = \{a\}$ and let $S = (\Sigma^*, \Sigma, <)$. We define

$$R := \mathrm{rep}_S(X) = \{(a^{n_1}, \ldots, a^{n_d}) \colon (n_1, \ldots, n_d) \in X\}.$$

The set X is 1-recognizable if and only if the language $R^\#$ is regular. By Theorem 2.3, the language $R^\#$ is regular if and only if it is a finite union of languages of the form

$$R_0 \ldots R_t, \quad t \in \mathbb{N},$$

where each factor $R_\ell \subseteq (\Sigma_{A_\ell})^*$ is regular and $A_t \subseteq \ldots \subseteq A_0 \subseteq \{1, \ldots, d\}$. Since $|\Sigma| = 1$, we have $|\Sigma_{A_\ell}| = 1$. Let $\Sigma_{A_\ell} = \{x\}$. It is well-known [6], Prop. V.1.1, that R_ℓ is a finite union of languages of the form $\{x^{pi+q} : i \in \mathbb{N}\}$, where $p, q \in \mathbb{N}$. Without loss of generality we can assume that R_ℓ is exactly of this form. Hence, the language R_ℓ consists of the representations of a set of the form

$$\{c_\ell(n_{\ell,1}, \ldots, n_{\ell,d}) + (b_{\ell,1}, \ldots, b_{\ell,d}) : (\forall i)(n_{\ell,i} \in \mathbb{N})\}.$$

The conditions $A_t \subseteq \ldots \subseteq A_0 \subseteq \{1, \ldots, d\}$ impose the restrictions on the $n_{\ell,i}$'s and the constants $b_{\ell,i}$ in the statement of the lemma. The concatenation of the R_ℓ's gives the sum described above. □

Remark 2.7. We can give an alternative description of the 1-recognizable sets. Let $\mathbf{v} = (v_1, \ldots, v_d) \in \mathbb{N}^d$. We define $\mathrm{Supp}(\mathbf{v}) = \{i \in \{1, \ldots, d\} : v_i \neq 0\}$. Let $X \subseteq \mathbb{N}^d$. Then X is a finite union of sets of the form described in Lemma 2.6 if and only if X is a finite union of sets of the form

$$(b_0 + c_0\,\mathbb{N})\mathbf{v_0} + \ldots + (b_t + c_t\,\mathbb{N})\mathbf{v_t},$$

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where

- $t \in \mathbb{N}$;
- $b_i, c_i \in \mathbb{N}$ for $i = 1, \ldots, t$;
- $\mathbf{v}_i \in \{0,1\}^d$ for $i = 1, \ldots, t$;
- $\mathrm{Supp}(\mathbf{v}_t) \subseteq \ldots \subseteq \mathrm{Supp}(\mathbf{v}_0)$.

Example 2.8. Let $X = \{(5n, 5n+4m+6\ell+1, 5n+4m+6\ell+3, 5n) : n, m, \ell \in \mathbb{N}\}$. The unary representation of X is

$$R^{\#} = ((a,a,a,a)^5)^*((\#,a,a,\#)^4)^*((\#,a,a,\#)^6)^*(\#,a,a,\#)(\#,\#,a,\#)^2.$$

Since $R^{\#}$ is regular the set X is 1-recognizable. The set X can be written as

$$X = \{5(n,n,n,n)+4(0,m,m,0)+6(0,\ell,\ell,0)+(0,1,1,0)+(0,0,2,0) : n, m, \ell \in \mathbb{N}\},$$

which is an expression of the form (2.1) where $t = 3$; $A_0 = \{1,2,3,4\}$, $A_1 = A_2 = \{2,3\}$, $A_3 = \{3\}$; $c_0 = 5, c_1 = 4, c_2 = 6, c_3 = 0$; and $b_{0,i} = b_{1,i} = 0$ for all i, $(b_{2,1}, b_{2,2}, b_{2,3}, b_{2,4}) = (0,1,1,0)$, $(b_{3,1}, b_{3,2}, b_{3,3}, b_{3,4}) = (0,0,2,0)$.

Alternatively, by Remark 2.7 we can write

$$X = 5\,\mathbb{N}(1,1,1,1) + 4\,\mathbb{N}(0,1,1,0) + (1 + 6\,\mathbb{N})(0,1,1,0) + (2 + 0\,\mathbb{N})(0,0,1,0).$$

Furthermore, we have a factorization of $R^{\#}$ as given in Theorem 2.3: that is, $R^{\#} = R_0R_1R_2R_3$, where $R_0 = ((a,a,a,a)^5)^*$, $R_1 = ((\#,a,a,\#)^4)^*$, $R_2 = ((\#,a,a,\#)^6)^*(\#,a,a,\#)$, and $R_3 = (\#,\#,a,\#)^2$, with the same A_ℓ's as those defined above. The term $5(n,n,n,n)$ corresponds to R_0, the term $4(0,m,m,0)$ corresponds to R_1, the term $6(0,\ell,\ell,0) + (0,1,1,0)$ corresponds to R_2, and the term $(0,0,2,0)$ corresponds to R_3.

We need the following classical number-theoretic result (see [13], Thm. 1.0.1).

Theorem 2.9. *Let $a_1, \ldots, a_n$ be integers with $a_i \geq 2$ for $i = 1, \ldots, n$. If*

$$\gcd(a_1, \ldots, a_n) = 1,$$

then there exists a positive integer $F(a_1, \ldots, a_n)$ such that $F(a_1, \ldots, a_n)$ cannot be expressed as a non-negative linear combination of $a_1, \ldots, a_n$, but all integers greater than $F(a_1, \ldots, a_n)$ can be so expressed.

In the sequel we write $\mathbf{e}_i$ to denote the element of $\mathbb{N}^d$ that contains a 1 in its i-th component and 0's in all others.

Lemma 2.10. *A set $X \subseteq \mathbb{N}^d$ of the form (2.1) can be written as a union $A \cup B$, where A is made up of finite unions and intersections of sets having one of the forms (2.2)–(2.5) below and B is a finite intersection of sets of the form (2.2) or (2.3) below:*

$$\left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (rn_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \geq N \right\} \tag{2.2}$$

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where $1 \le j \le d$, *and* $r, s, N \in \mathbb{N}$;

$$\left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (n_k + r n_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \ge N \right\} \tag{2.3}$$

where $1 \le j, k \le d$, $j \neq k$, *and* $r, s, N \in \mathbb{N}$;

$$\left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (r n_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \in C \right\} \tag{2.4}$$

where $1 \le j \le d$, $r, s \in \mathbb{N}$, *and* $C \subseteq \mathbb{N}$ *is a finite set; or*

$$\left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (n_k + r n_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \in C \right\} \tag{2.5}$$

where $1 \le j, k \le d$, $j \neq k$, *and* $r, s \in \mathbb{N}$, *and* $C \subseteq \mathbb{N}$ *is a finite set.*

Proof. Let X be a set of the form (2.1) where t, the A_ℓ's, the c_ℓ's, and the $b_{\ell,i}$'s are fixed and satisfy the conditions listed in Lemma 2.6. We will write $X = A \cup B$, where

$$B = \bigcap_{j=1}^{d} Y_j,$$

where each Y_j is either of the form (2.2) or (2.3), and A is made up of finite unions and intersections of sets of the forms (2.2)–(2.5).

First observe that if $j \in \{1, \ldots, d\} \setminus A_0$ the set X contains only vectors whose j-th component is always 0. For each such j, we define

$$Y_j = \left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + 0\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N} \right\},$$

which is of the form (2.2).

First consider the case where $A_0 = \ldots = A_t$. Define $j_1 < \ldots < j_{|A_0|}$ to be the elements of A_0. Define

$$Y_{j_1} = \left\{ \sum_{\substack{i=1 \\ i \neq j_1}}^{d} n_i \mathbf{e}_i + (r n_{j_1} + s)\mathbf{e}_{j_1} : n_1, \ldots, n_d \in \mathbb{N}, n_{j_1} \ge N \right\},$$

where $r = \gcd(c_0, \ldots, c_t)$, $s = \sum_{\ell=0}^{t} b_{\ell,j_1}$, and $N - 1$ is the largest integer n such that rn cannot be written as a nonnegative integer linear combination of

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$c_0, \ldots, c_t$ (note that by Theorem 2.9, N exists and is finite). Note that Y_{j_1} is of the form (2.2).

Define

$$Y'_{j_1} = \left\{ \sum_{\substack{i=1 \\ i \neq j_1}}^{d} n_i \mathbf{e}_i + (rn_{j_1} + s)\mathbf{e}_{j_1} : n_1, \ldots, n_d \in \mathbb{N}, n_{j_1} \in C \right\},$$

where C is the set of all nonnegative integers $n < N$ such that rn can be written as a nonnegative integer linear combination of $c_0, \ldots, c_t$. Note that Y'_{j_1} is of the form (2.4).

For $k \in \{2, \ldots, |A_0|\}$, define

$$Y_{j_k} = \left\{ \sum_{\substack{i=1 \\ i \neq j_k}}^{d} n_i \mathbf{e}_i + n_{j_{k-1}} \mathbf{e}_{j_k} : n_1, \ldots, n_d \in \mathbb{N} \right\},$$

which is of the form (2.3).

The set X can be written as the union $A \cup B$ where

$$B = \bigcap_{j \in \{1,\ldots,d\} \setminus A_0} Y_j \quad \cap \quad \bigcap_{k \in \{1,\ldots,|A_0|\}} Y_{j_k}$$

and

$$A = \bigcap_{j \in \{1,\ldots,d\} \setminus A_0} Y_j \quad \cap \quad \bigcap_{k \in \{2,\ldots,|A_0|\}} Y_{j_k} \quad \cap \quad Y'_{j_1}.$$

Now consider the case where there is at least one index ℓ such that $A_\ell \setminus A_{\ell+1} \neq \emptyset$. Define $\ell_1 < \ldots < \ell_{t'}$ to be the indices of the sets A_ℓ satisfying $A_{\ell_k} \setminus A_{\ell_k+1} \neq \emptyset$ for each $k \in \{1, \ldots, t'\}$. We clearly have $1 \leq t' \leq t$ and $0 \leq \ell_{t'} < t$.

Define $d_1 = |A_{\ell_1} \setminus A_{\ell_1+1}|$ and $j_{1,1} < \ldots < j_{1,d_1}$ to be the elements of $A_{\ell_1} \setminus A_{\ell_1+1}$. Define

$$Y_{j_{1,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{1,1}}}^{d} n_i \mathbf{e}_i + (r_1 n_{j_{1,1}} + s_1)\mathbf{e}_{j_{1,1}} : n_1, \ldots, n_d \in \mathbb{N}, n_{j_{1,1}} \geq N_1 \right\},$$

where $r_1 = \gcd(c_0, \ldots, c_{\ell_1})$, $s_1 = \sum_{\ell=0}^{\ell_1} b_{\ell, j_{1,1}}$, and $N_1 - 1$ is the largest integer n such that $r_1 n$ cannot be written as a nonnegative integer linear combination of $c_0, \ldots, c_{\ell_1}$. Note that $Y_{j_{1,1}}$ is of the form (2.2).

Define

$$Y'_{j_{1,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{1,1}}}^{d} n_i \mathbf{e}_i + (r_1 n_{j_{1,1}} + s_1)\mathbf{e}_{j_{1,1}} : n_1, \ldots, n_d \in \mathbb{N}, n_{j_{1,1}} \in C_1 \right\},$$

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where C_1 is the set of all nonnegative integers $n < N_1$ such that $r_1 n$ can be written as a nonnegative integer linear combination of $c_0, \dots, c_{\ell_1}$. Note that $Y'_{j_{1,1}}$ is of the form (2.4).

For $k \in \{2, \dots, d_1\}$, define

$$Y_{j_{1,k}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{1,k}}}^{d} n_i \mathbf{e}_i + n_{j_{1,k-1}} \mathbf{e}_{j_{1,k}} : n_1, \dots, n_d \in \mathbb{N} \right\},$$

which is of the form (2.3).

Define $d_2 = |A_{\ell_2} \setminus A_{\ell_2+1}|$ and $j_{2,1} < \ldots < j_{2,d_2}$ to be the elements of $A_{\ell_2} \setminus A_{\ell_2+1}$. Define

$$Y_{j_{2,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{2,1}}}^{d} n_i \mathbf{e}_i + (n_{j_{1,1}} + r_2 n_{j_{2,1}} + s_2) \mathbf{e}_{j_{2,1}} : n_1, \dots, n_d \in \mathbb{N}, n_{j_{2,1}} \geq N_2 \right\},$$

where $r_2 = \gcd(c_{\ell_1+1}, \dots, c_{\ell_2})$, $s_2 = \sum_{\ell=\ell_1+1}^{\ell_2} b_{\ell,j_{2,1}}$, and $N_2 - 1$ is the largest integer n such that $r_2 n$ cannot be written as a nonnegative integer linear combination of $c_{\ell_1+1}, \dots, c_{\ell_2}$. Note that $Y_{j_{2,1}}$ is of the form (2.3).

Define

$$Y'_{j_{2,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{2,1}}}^{d} n_i \mathbf{e}_i + (n_{j_{1,1}} + r_2 n_{j_{2,1}} + s_2) \mathbf{e}_{j_{2,1}} : n_1, \dots, n_d \in \mathbb{N}, n_{j_{2,1}} \in C_2 \right\},$$

where C_2 is the set of all nonnegative integers $n < N_2$ such that $r_2 n$ can be written as a nonnegative integer linear combination of $c_{\ell_1+1}, \dots, c_{\ell_2}$. Note that $Y'_{j_{2,1}}$ is of the form (2.5).

For $k \in \{2, \dots, d_2\}$, define

$$Y_{j_{2,k}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{2,k}}}^{d} n_i \mathbf{e}_i + n_{j_{2,k-1}} \mathbf{e}_{j_{2,k}} : n_1, \dots, n_d \in \mathbb{N} \right\},$$

which is of the form (2.3).

We continue in this manner to define d_p, $Y_{j_{p,k}}$, and $Y'_{j_{p,1}}$ for all $p \in \{1, \dots, t'\}$ and $k \in \{1, \dots, d_p\}$. Finally observe that we have $A_{\ell_{t'}} \setminus A_{\ell_{t'}+1} \neq \emptyset$ and $A_{\ell_{t'}+1} = \ldots = A_t$. Define $d_{t'+1} = |A_t|$ and $j_{t'+1,1} < \ldots < j_{t'+1,d_{t'+1}}$ to be the elements of

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A_t. Define

$$Y_{j_{t'+1,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{t'+1,1}}}^{d} n_i \mathbf{e}_i + (n_{j_{t',1}} + r_{t'+1} n_{j_{t'+1,1}} + s_{t'+1}) \mathbf{e}_{j_{t'+1,1}} : \right.$$
$$\left. n_1, \ldots, n_d \in \mathbb{N}, n_{j_{t'+1,1}} \geq N_{t'+1} \right\},$$

where $r_{t'+1} = \gcd(c_{\ell_{t'}+1}, \ldots, c_t)$, $s_{t'+1} = \sum_{\ell=\ell_{t'}+1}^{\ell_t} b_{\ell, j_{t'+1,1}}$, and $N_{t'+1} - 1$ is the largest integer n such that $r_{t'+1} n$ cannot be written as a nonnegative integer linear combination of $c_{\ell_{t'}+1}, \ldots, c_t$. Again note that $Y_{j_{t'+1,1}}$ is of the form (2.3).

Define

$$Y'_{j_{t'+1,1}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{t'+1,1}}}^{d} n_i \mathbf{e}_i + (n_{j_{t',1}} + r_{t'+1} n_{j_{t'+1,1}} + s_{t'+1}) \mathbf{e}_{j_{t'+1,1}} : \right.$$
$$\left. n_1, \ldots, n_d \in \mathbb{N}, n_{j_{t'+1,1}} \in C_{t'+1} \right\},$$

where $C_{t'+1}$ is the set of all nonnegative integers $n < N_{t'+1}$ such that $r_{t'+1} n$ can be written as a nonnegative integer linear combination of $c_{\ell_{t'}+1}, \ldots, c_t$. Note that $Y'_{j_{t'+1,1}}$ is of the form (2.5).

For $k \in \{2, \ldots, d_{t'+1}\}$, define

$$Y_{j_{t'+1,k}} = \left\{ \sum_{\substack{i=1 \\ i \neq j_{t'+1,k}}}^{d} n_i \mathbf{e}_i + n_{j_{t'+1,k-1}} \mathbf{e}_{j_{t'+1,k}} : n_1, \ldots, n_d \in \mathbb{N} \right\},$$

which is of the form (2.3).

The set X can be written as the union $A \cup B$ where

$$B = \bigcap_{j \in \{1,\ldots,d\} \setminus A_0} Y_j \quad \cap \quad \bigcap_{\substack{p \in \{1,\ldots,t'+1\} \\ k \in \{1,\ldots,d_p\}}} Y_{j_{p,k}}$$

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and

$$A = \bigcap_{j\in\{1,\dots,d\}\setminus A_0} Y_j \quad \cap \quad \bigcup_{p\in\{1,\dots,t'+1\}} \left(Y'_{j_{p,1}} \quad \cap \quad \bigcap_{\substack{q\in\{1,\dots,t'+1\}\setminus\{p\}\\ k\in\{1,\dots,d_q\}}} Y_{j_{q,k}} \quad \cap \quad \bigcap_{k\in\{2,\dots,d_p\}} Y_{j_{p,k}} \right). \quad \square$$

Example 2.11. We continue Example 2.8. We will write $X = A \cup B$ as in Lemma 2.10. The A_ℓ's are not all the same, so we can define $t' = 2$, $\ell_1 = 0 < \ell_2 = 2$ as in the proof of Lemma 2.10.

We have $d_1 = |A_0 \setminus A_1| = 2$, $j_{1,1} = 1$ and $j_{1,2} = 4$. We also have $r_1 = \gcd(c_0) = \gcd(5) = 5$ and $s_1 = 0$, and hence $N_1 = 0$. Therefore,

$$Y_1 = \{n_2\mathbf{e_2} + n_3\mathbf{e_3} + n_4\mathbf{e_4} + (5n_1 + 0)\mathbf{e_1} : n_1, n_2, n_3, n_4 \in \mathbb{N}, n_1 \geq 0\},$$

$$Y'_1 = \{n_2\mathbf{e_2} + n_3\mathbf{e_3} + n_4\mathbf{e_4} + (5n_1 + 0)\mathbf{e_1} : n_2, n_3, n_4 \in \mathbb{N}, n_1 \in C_1\} = \emptyset,$$

since $C_1 = \emptyset$, and

$$Y_4 = \{n_1\mathbf{e_1} + n_2\mathbf{e_2} + n_3\mathbf{e_3} + n_1\mathbf{e_4} : n_1, n_2, n_3 \in \mathbb{N}\}.$$

Next we have $d_2 = |A_2 \setminus A_3| = 1$ and $j_{2,1} = 2$. We also have $r_2 = \gcd(c_1, c_2) = \gcd(4, 6) = 2$ and $s_2 = b_{1,2} + b_{2,2} = 0 + 1 = 1$, and hence $N_2 = 2$. Therefore,

$$Y_2 = \{n_1\mathbf{e_1} + n_3\mathbf{e_3} + n_4\mathbf{e_4} + (n_1 + 2n_2 + 1)\mathbf{e_2} : n_1, n_2, n_3 \in \mathbb{N}, n_2 \geq 2\},$$

and

$$\begin{aligned} Y'_2 &= \{n_1\mathbf{e_1} + n_3\mathbf{e_3} + n_4\mathbf{e_4} + (n_1 + 2n_2 + 1)\mathbf{e_2} : n_1, n_3 \in \mathbb{N}, n_2 \in C_2\} \\ &= \{n_1\mathbf{e_1} + n_3\mathbf{e_3} + n_4\mathbf{e_4} + (n_1 + 1)\mathbf{e_2} : n_1, n_3 \in \mathbb{N}\}, \end{aligned}$$

since $C_2 = \{0\}$.

Finally, we have $d_3 = |A_3| = 1$ and $j_{3,1} = 3$. We also have $r_3 = \gcd(c_3) = \gcd(0) = 0$ and $s_3 = b_{3,3} = 2$, and hence $N_3 = 0$. Therefore,

$$Y_3 = \{n_1\mathbf{e_1} + n_2\mathbf{e_2} + n_4\mathbf{e_4} + (n_2 + 0n_3 + 2)\mathbf{e_3} : n_1, n_2, n_3 \in \mathbb{N}, n_3 \geq 0\},$$

and

$$Y'_3 = \{n_1\mathbf{e_1} + n_2\mathbf{e_2} + n_4\mathbf{e_4} + (n_2 + 0n_3 + 2)\mathbf{e_3} : n_1, n_2 \in \mathbb{N}, n_3 \in C_3\} = \emptyset,$$

since $C_3 = \emptyset$.

Hence $A = Y_1 \cap Y'_2 \cap Y_3 \cap Y_4$ and $B = Y_1 \cap Y_2 \cap Y_3 \cap Y_4$.

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Lemma 2.12. *Let $k \in \mathbb{N}$ and let S be an abstract numeration system. The set $X = \{(n, n+k) : n \in \mathbb{N}\}$ is S-recognizable.*

Proof. The proof follows easily from known results and so we only give a sketch of the proof. Let $R = \mathrm{rep}_S(X)$. To show that X is S-recognizable we must show that $R^{\#}$ is a regular language. Consider first the set $Y = \{(\mathrm{rep}_S(n), \mathrm{rep}_S(n+1)) : n \in \mathbb{N}\}$. If we interpret Y as the function mapping $\mathrm{rep}_S(n)$ to $\mathrm{rep}_S(n+1)$, then Y is the so-called *successor function* (see [1] or [11] for more on the successor function). From [2], Proposition 3 (see also [9], Prop. 2.6.7), we have that Y is a *synchronous relation*. In [9] synchronous relations are defined in terms of letter-to-letter transducers, but this definition is equivalent to the fact that the language $Y^{\#}$ is accepted by a finite automaton. Moreover, from [8] (see also [9], Thm. 2.6.6), we have that the composition of synchronous relations is again a synchronous relation. Hence R, which is the k-fold composition of Y with itself, is a synchronous relation. We conclude that $R^{\#}$ is a regular language, as required. □

Lemma 2.13. *A set $X \subseteq \mathbb{N}^d$ having one of the forms (2.2)–(2.5) defined in Lemma 2.10 is S-recognizable for any abstract numeration system S.*

Proof. We will give the proof for the cases where X is either of the form (2.2) or (2.3) (the other two cases are similar).

Let $S = (L, \Sigma, <)$ be an abstract numeration system and let $\mathcal{T}$ be a finite automaton accepting L. Let $R = \mathrm{rep}_S(X)$. We will show that $R^{\#}$ is regular. That is, we will define a (nondeterministic) finite automaton $\mathcal{M}$ that accepts $R^{\#}$. Let $(w_1, \ldots, w_d)^{\#}$ be an arbitrary input to the automaton $\mathcal{M}$.

Suppose that X is of the form (2.2). That is,

$$X = \left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (rn_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \geq N \right\},$$

where $1 \leq j \leq d$, and $r, s, N \in \mathbb{N}$. Suppose first that $r = 0$. In this case, the automaton $\mathcal{M}$ simulates $\mathcal{T}$ on $w_1, \ldots, w_{j-1}, w_{j+1}, \ldots, w_d$. The automaton $\mathcal{M}$ accepts its input if and only if $\mathcal{T}$ accepts $w_1, \ldots, w_{j-1}, w_{j+1}, \ldots, w_d$ and $w_j = \mathrm{rep}_S(s)$.

Now suppose that $r > 0$. By increasing the value of N, we may, without loss of generality, assume that $s < r$. By Theorem 1.3 the language $\{\mathrm{rep}_S(rn_j + s) : n_j \geq N\}$ is regular. Let $\mathcal{T}'$ be an automaton accepting $\{\mathrm{rep}_S(rn_j + s) : n_j \geq N\}$. As before, the automaton $\mathcal{M}$ simulates $\mathcal{T}$ on $w_1, \ldots, w_{j-1}, w_{j+1}, \ldots, w_d$, but now also simulates $\mathcal{T}'$ on w_j. The automaton $\mathcal{M}$ accepts its input if and only if $\mathcal{T}$ accepts $w_1, \ldots, w_{j-1}, w_{j+1}, \ldots, w_d$ and $\mathcal{T}'$ accepts w_j.

Next suppose that X is of the form (2.3). That is,

$$\left\{ \sum_{\substack{i=1 \\ i \neq j}}^{d} n_i \mathbf{e}_i + (n_k + rn_j + s)\mathbf{e}_j : n_1, \ldots, n_d \in \mathbb{N}, n_j \geq N \right\},$$

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where $1 \leq j, k \leq d$, $j \neq k$, and $r, s, N \in \mathbb{N}$. Again, suppose first that $r = 0$. By Lemma 2.12, the language $\{(\mathrm{rep}_S(n_k), \mathrm{rep}_S(n_k + s))^\# : n_k \in \mathbb{N}\}$ is regular. Let $\mathcal{T}''$ be a finite automaton accepting this language. The automaton $\mathcal{M}$ simulates $\mathcal{T}$ on each of the words in $\{w_1, \ldots, w_d\} \setminus \{w_j, w_k\}$. Simultaneously, the automaton $\mathcal{M}$ simulates $\mathcal{T}''$ on the pair $(w_k, w_j)^\#$. The automaton $\mathcal{M}$ accepts its input if and only if $\mathcal{T}$ accepts $\{w_1, \ldots, w_d\} \setminus \{w_j, w_k\}$ and $\mathcal{T}''$ accepts $(w_k, w_j)^\#$.

Now suppose that $r > 0$. Again, without loss of generality, we may assume that $s < r$. Using the same ideas as in the proof of [12], Theorem 3.3.1, it is not hard to see that the language

$$\{(\mathrm{rep}_S(m), \mathrm{rep}_S(n))^\# : m, n \in \mathbb{N} \text{ and } (n - m) \equiv s \pmod{r}\}$$

is regular. Let $\mathcal{Z}$ be an automaton accepting this language. Let $\mathcal{Z}'$ be an automaton accepting the language $\{(\mathrm{rep}_S(n_k), \mathrm{rep}_S(n_k + rN + s))^\# : n_k \in \mathbb{N}\}$ (since $rN + s$ is a constant, we may apply Lem. 2.12).

The automaton $\mathcal{M}$ simulates $\mathcal{T}$ on each of the words in $\{w_1, \ldots, w_d\} \setminus \{w_j, w_k\}$. Simultaneously, the automaton $\mathcal{M}$ simulates $\mathcal{Z}$ on the pair $(w_k, w_j)^\#$.

The automaton $\mathcal{M}$ also nondeterministically "guesses" a word $v = b_1 \ldots b_{|v|}$ and simulates $\mathcal{Z}'$ on the pair $(w_k, v)^\#$. This "guess" works as follows. Let $w_k = a_1 \ldots a_{|w_k|}$, where each $a_i \in \Sigma$. For each $i = 1, \ldots, |w_k|$, we simulate $\mathcal{Z}'$ by nondeterministically choosing to follow one of the transitions of $\mathcal{Z}'$ labeled (a_i, b_i), where $b_i \in \Sigma$; and for $i > |w_k|$ (*i.e.*, w_k has been completely read), the simulation may make a nondeterministic choice among transitions of the form $(\#, b_i)$, where $b_i \in \Sigma$. This nondeterministic choice of b_i at each step of the simulation is what defines the "guessed" word v. Note that if $\mathcal{Z}'$ accepts $(w_k, v)^\#$, then $\mathrm{val}_s(v) = \mathrm{val}_S(w_k) + rN + s$. As this nondeterministic simulation is performed, the automaton $\mathcal{M}$ also simultaneously verifies that w_j is greater than or equal to (in the radix order) the guessed word v.

The automaton $\mathcal{M}$ accepts its input if and only if

- $\mathcal{T}$ accepts each of the words in $\{w_1, \ldots, w_d\} \setminus \{w_j, w_k\}$;
- $\mathcal{Z}$ accepts $(w_k, w_j)^\#$ (and hence $\mathrm{val}(w_j) - \mathrm{val}(w_k) \equiv s \pmod{r}$);
- $\mathcal{Z}'$ accepts $(w_k, v)^\#$ for some guessed word v as described above (and hence $\mathrm{val}_S(v) = \mathrm{val}_S(w_k) + rN + s$); and
- w_j is greater than or equal to v in the radix order (and hence $\mathrm{val}_S(w_j) \geq \mathrm{val}_S(v)$).

The last three of these conditions guarantee that $\mathrm{val}_S(w_j) = \mathrm{val}_S(w_k) + rn_j + s$ for some $n_j \geq N$.

This completes the proof for the cases where X is either of the form (2.2) or (2.3). As previously stated, we omit the details for the other two cases since they are similar. □

We are ready for the proof of Theorem 1.9.

Proof of Theorem 1.9. One direction is clear: if X is S-recognizable for all abstract numeration systems S, then it is certainly 1-recognizable.

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To prove the other direction, suppose that X is 1-recognizable. The result now follows from Lemmas 2.6, 2.10, and 2.13. □

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